## Computability and Logic

## HW 7

## Due: Friday, April 17

1. Use the Abacus machine software to create an Abacus machine that computes max $(x, y)$.
2. Use the Abacus machine software to create an Abacus machine that computes the quo( $\mathrm{x}, \mathrm{y}$ ) and rem $(x, y)$ functions at the same time by leaving quo in register 3 , and rem in register 4 (make sure to leave x in register $1, \mathrm{y}$ in register 2 when the machine halts, and empty out all registers beyond 4). For the special case where $y=0$ we set quo $(x, 0)=0$ and rem $(x, 0)=x$.
3. Questions A, B, C from the Abacus-Machine Halting Function handout.
4. Use the Turing-machine software to create a Turing machine that simulates the following Abacus machine:


Your Turing-machine should simulate this Abacus-machine using the following simulation conventions:

- Assume that the Abacus-machine has more than 2 registers (i.e. act as if this Abacus-machine is part of a larger machine)
- The Turing-tape is used to represent the contents of the registers using unary representation. To be specific, if the Abacus-machine registers $R_{1}, R_{2}, \ldots R_{k}$ contain the numbers $n_{1}, n_{2}, \ldots n_{k}$ respectively, then this will be represented using Turing-Tape configuration $\left[n_{1}, n_{2}, \ldots n_{k}\right]$ (as defined in HW 6)
- The Turing-machine starts at the leftmost 1 on the tape, and after each instruction of the Abacus-machine that the Turing-machine simulates, the Turing-machine should return the head back to the leftmost one.

So I am *not* looking for a Turing-machine that accomplishes whatever this Abacus-machine accomplishes merely in terms input and output behavior, but for a Turing-machine that does it using the above conventions!

For all Abacus-machines you create for the HW, follow the following guidelines:

- Use the Abacus machine software AMS 2.1 to create your machine and submit electronically.
- Unless stated otherwise, when using Abacus-machines to compute functions over natural numbers, use the 'standard' convention of changing input configuration $\left[n_{1}, n_{2}, \ldots n_{k}\right]$ into output configuration $\left[n_{1}, n_{2}, \ldots n_{k} f\left(n_{1}, n_{2}, \ldots n_{k}\right)\right]$ where $\left[n_{1}, n_{2}, \ldots n_{k}\right]$ represents a tuple of numbers $<n_{1}$, $n_{2}, \ldots n_{k}>$ as follows: for all $1 \leq i \leq n$ : register $R_{i}$ contains $n_{i}$, and all other registers are empty. In other words: make sure that when the machine is done, you have the input of $k$ numbers back in the first $k$ registers, the function value is in register $k+1$, and all other registers are empty.
- Organize the nodes and transitions so that your machine looks nice and is 'readable' (i.e. don't leave your machine a spaghetti of connections). Try to keep crossing connections at a minimum.

